## Exercise 7.1.

Trace the TM in Figure 7.6 (see next slide), accepting the language $\left\{x x \mid x \in\{a, b\}^{*}\right\}$, on the string aaba. Show the configuration at each step.


## Exercise 7.2.

Below is a transition table for a TM with input alphabet $\{a, b\}$.

| $q$ | $\sigma$ | $\delta(q, \sigma)$ | $q$ | $\sigma$ | $\delta(q, \sigma)$ | $q$ | $\sigma$ | $\delta(q, \sigma)$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $q_{0}$ | $\Delta$ | $\left(q_{1}, \Delta, R\right)$ | $q_{2}$ | $\Delta$ | $\left(h_{a}, \Delta, R\right)$ | $q_{6}$ | $a$ | $\left(q_{6}, a, R\right)$ |
| $q_{1}$ | $a$ | $\left(q_{1}, a, R\right)$ | $q_{3}$ | $\Delta$ | $\left(q_{4}, a, R\right)$ | $q_{6}$ | $b$ | $\left(q_{6}, b, R\right)$ |
| $q_{1}$ | $b$ | $\left(q_{1}, b, R\right)$ | $q_{4}$ | $a$ | $\left(q_{4}, a, R\right)$ | $q_{6}$ | $\Delta$ | $\left(q_{7}, b, L\right)$ |
| $q_{1}$ | $\Delta$ | $\left(q_{2}, \Delta, L\right)$ | $q_{4}$ | $b$ | $\left(q_{4}, b, R\right)$ | $q_{7}$ | $a$ | $\left(q_{7}, a, L\right)$ |
| $q_{2}$ | $a$ | $\left(q_{3}, \Delta, R\right)$ | $q_{4}$ | $\Delta$ | $\left(q_{7}, a, L\right)$ | $q_{7}$ | $b$ | $\left(q_{7}, b, L\right)$ |
| $q_{2}$ | $b$ | $\left(q_{5}, \Delta, R\right)$ | $q_{5}$ | $\Delta$ | $\left(q_{6}, b, R\right)$ | $q_{7}$ | $\Delta$ | $\left(q_{2}, \Delta, L\right)$ |

What is the final configuration if the TM starts with input string $x$ ?

## Exercise 7.3.

Let $T=\left(Q, \Sigma, \Gamma, q_{0}, \delta\right)$ be a TM, and let $s$ and $t$ be the sizes of the sets $Q$ and $\Gamma$, respectively.
How many distinct configurations of $T$ could there possibly be in which all tape squares past square $n$ are blank and $T$ 's tape head is on or to the left of square $n$ ? (The tape squares are numbered beginning with 0 .)

## Exercise 7.10.

We do not define $\wedge$-transitions for a TM. Why not? What features of a TM make it unnecessary or inappropriate to talk about $\wedge$-transitions?

## Exercise 7.17.

For each case below, draw a TM that computes the indicated function.

In the first four parts, the function is from $\mathbb{N}$ to $\mathbb{N}$. In each of these parts, assume that the TM uses unary notation - i.e., the natural number $n$ is represented by the string $1^{n}$.
a. $f(x)=x+2$
b. $f(x)=2 x$
c. $f(x)=x^{2}$
e. $E:\{a, b\}^{*} \times\{a, b\}^{*} \rightarrow\{0,1\}$ defined by $E(x, y)=1$ if $x=y, \quad E(x, y)=0$ otherwise.

## Exercise.

Draw a TM that computes the function

$$
f(x, y)=x+y
$$

where $x, y$ are integers $\geq 0$.

Assume that the TM uses unary notation, both for its input and for its output.

Make this exercise yourself.

## Exercise.

Draw a TM that computes the function $f(x, y)=x \bmod y$

Hint: implement the following algorithm:

$$
\begin{gathered}
\text { while }(x>=y) \\
x=x-y
\end{gathered}
$$

Make this exercise yourself.

## Exercise 7.12.

Suppose $T$ is a TM that accepts a language $L$.
Describe how you would modify $T$ to obtain another TM accepting $L$ that never halts in the reject state $h_{r}$.

## Exercise 7.16.

Does every TM compute a partial function? Explain.

